Importance Function Derivation for RESTART Simulations of Petri Nets

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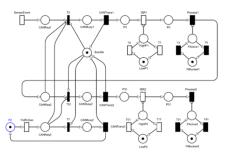
Outline

- Stochastic Petri Nets
- RESTART Simulation
- A-Priori Estimation of Rare Event Frequencies
- A Heuristic Importance Function for SPNs
- An Example
- Conclusion

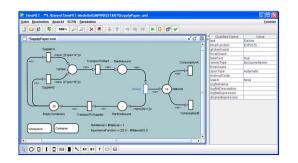
Introduction

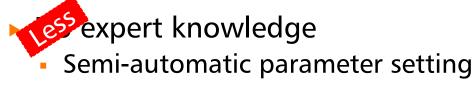
Accessible Rare-Event Simulation

- Underlying model
 - Stochastic Petri nets



- Software Tool
 - TimeNET



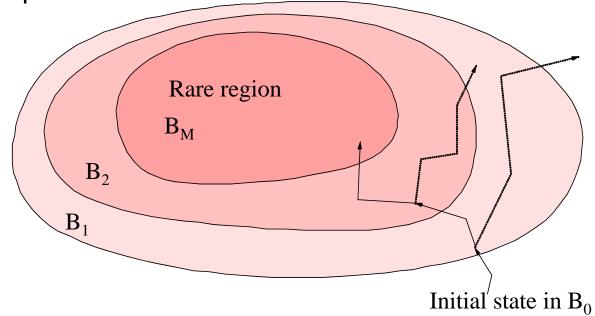




RESTART

Splitting Simulation

- Model and simulation loop as (semi-) black box
- Controlling RESTART algorithm only needs to …
 - Store and recall state
 - Decide when to split and how often
 - Technical: Importance function and Levels
 - Crucial for speedup!



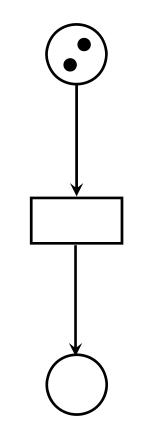
Stochastic Petri Nets

• Model Elements

- ► Places *P* state variables
 - Marking state, initial: \mathbf{m}_0
- ► Transitions T events
 - Input arcs $\operatorname{\mathbf{Pre}}: P \times T \to \mathbb{N}$
 - Arc multiplicities
 - Enabling and firing

Performance measures

- Rate and impulse rewards
- Concentrate on rare transition t_{rare} firing frequency γ





Stochastic Petri Nets

Algebraic Description

- Marking vector $\mathbf{m} \in \mathbb{N}^{|P|}$ and direct reachability $\mathbf{m} \stackrel{t}{\longrightarrow} \mathbf{m}'$
- Incidence matrix C with dimensions $|P| \times |T|$ C(p,t) denotes token change in p when t is fired
- ▶ Next marking $m' = m + C \cdot (0, ..., 0, 1_t, 0, ..., 0)$
- State equation $\mathbf{m} = \mathbf{m}_0 + \mathbf{C} \sigma$ for a firing count vector $\sigma \in \mathbb{N}^{|T|}$
- ► Invariants (P- and T-semiflows) $\mathbf{y} \mathbf{C} = 0 \longrightarrow \mathbf{y} \mathbf{m} = \mathbf{y} \mathbf{m}_0 + \mathbf{y} \mathbf{C} \sigma = \mathbf{y} \mathbf{m}_0 = \text{const.}$ $\mathbf{C} \mathbf{x} = 0 \longrightarrow \mathbf{m}' = \mathbf{m} + \mathbf{C} \mathbf{x} = \mathbf{m}$

Estimation of Performance Measures

- [Campos et. al.]
- Structural analysis of model
 constraints for upper and lower bounds of results
- Based on incidence matrix and other model information
- Formulation as a linear programming problem
- Notation: firing frequency = throughput
 inverse of average interfiring time

$$\gamma = \chi[t_{rare}] = \frac{1}{\Gamma[t_{rare}]}$$



- Visit Ratios and Service Demands
 - Derive relative visit ratios $\mathbf{v}^{(1)}: T \to \mathbb{R}^+$ of transitions using incidence matrix, transition invariants, and firing probabilities of conflicting transitions, relative to t_{rare}
 - ▶ Define average service demands D⁽¹⁾: T → ℝ⁺ relative to a transition by multiplying visit ratios by a transition's mean service time s̄[t_i] (its delay)

$$\overline{\mathbf{D}}^{(1)}[t_i] = \mathbf{v}^{(1)}[t_i]\,\overline{s}[t_i]$$



LPP Formulation for Bounds

Apply Little's Law to each (timed) transition and its input places for (m denotes the mean marking vector)

$$\Gamma[t_{rare}] \,\overline{\mathbf{m}} \geq \mathbf{Pre} \cdot \overline{\mathbf{D}}^{(1)}$$

- ▶ and thus $\Gamma[t_1] \ge \max_{\mathbf{y} \in \{\text{P-semiflows}\}} \frac{\mathbf{y} \cdot \mathbf{Pre} \cdot \overline{\mathbf{D}}^{(1)}}{\mathbf{y} \cdot \mathbf{m}_0}$
- Leading to a bound for the mean interfiring time

$$\Gamma[t_{rare}] = \begin{array}{ccc} \text{maximum} & \mathbf{y} \cdot \mathbf{Pre} \cdot \overline{\mathbf{D}}^{(1)} \\ \text{subject to} & \mathbf{y} \cdot \mathbf{C} &= 0 \\ & \mathbf{y} \cdot \mathbf{m}_0 &= 1 \\ & \mathbf{y} &\geq 0 \end{array}$$



- Throughput Estimation
 - ► Throughput upper bound $\chi_+[t_{rare}] = \frac{1}{\Gamma[t_{rare}]}$
 - Throughput lower bounds: worst-case interfiring time after all other transitions have fired using visit ratios

$$\chi_{-}[t_{1}] = \frac{1}{\sum_{t \in T} \overline{\mathbf{D}}^{(1)}[t]}$$
$$\chi_{-}[t_{j}] = \chi_{-}[t_{1}] \mathbf{v}^{(1)}[t_{j}]$$

Heuristic estimate: weighted average of bounds

$$\gamma = \chi[t_{rare}] \approx \alpha \chi_{-}[t_{rare}] + (1 - \alpha) \chi_{+}[t_{rare}]$$

Heuristic Importance Function

State Distances in Petri Nets

Following similar ideas, the minimal number of events (transition firings) h(m, t_{rare}) to reach a marking m' enabling transition t_{rare} is

- ► Admissible and monotone ⇒ useful for search problems
- Can be generalized to general marking conditions

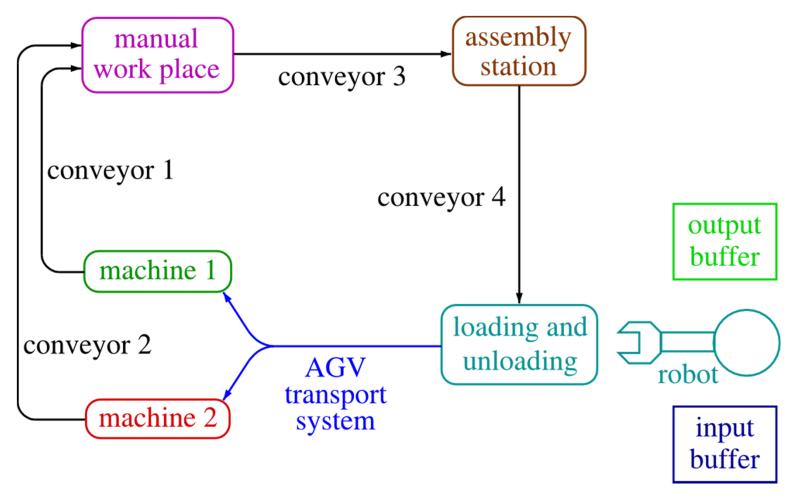
Heuristic Importance Function

- Importance Function Definition
 - Let $h_{\max} = h(\mathbf{m}_0, t_{rare})$, the maximum^{*} state distance
 - Importance function for marking m:

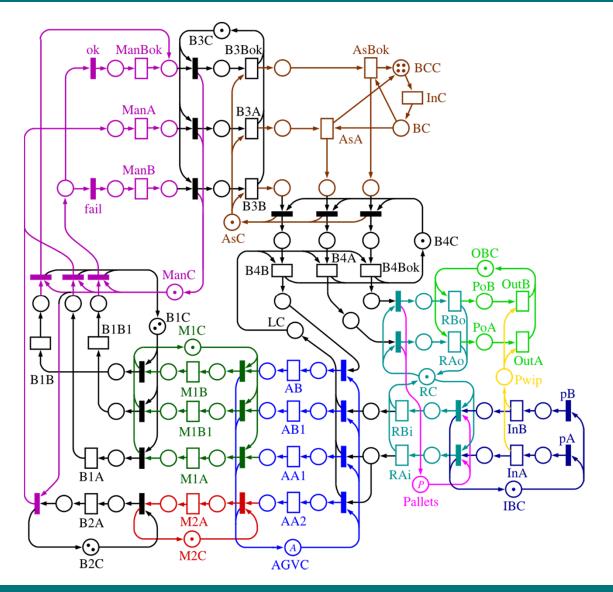
$$f_I(\mathbf{m}) = h_{\max} - h(\mathbf{m}, t_{rare})$$



Flexible Manufacturing System

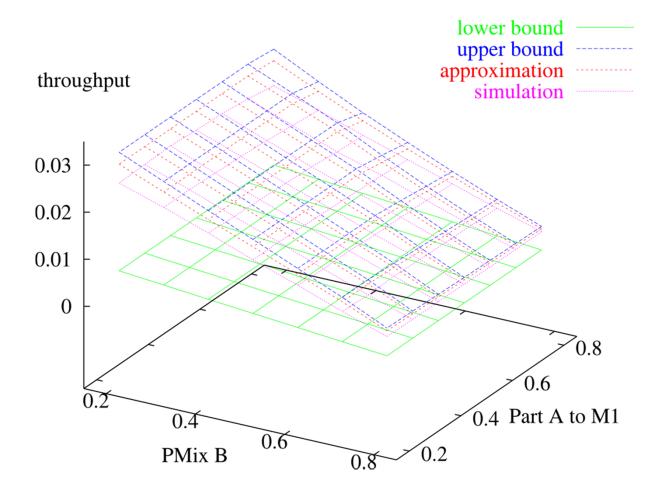


- Stochastic
 Petri Net
 (GSPN)
 - Different parts in various states





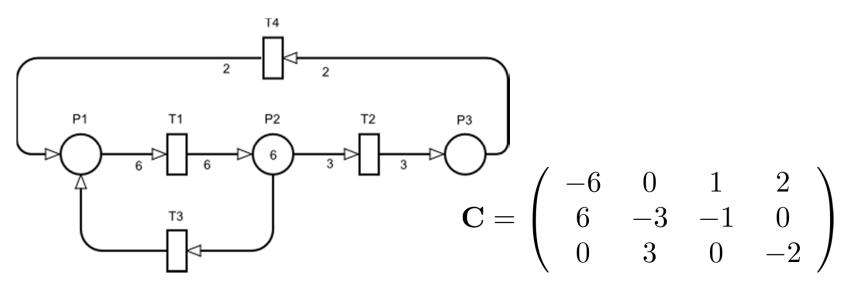
Throughput Estimation Quality





State Distance Heuristic

Simple model (manual solution)



How many transition firings until T1 is enabled?
 5: 2x T2, 3x T4

 (if arc multiplicities of T2 or T4 are decreased, then 6x T3)

Conclusion

Improvement of RESTART usability for SPNs

- A-priori estimation of rare event frequency
 Optimal number of thresholds
- Heuristic state distance measure
 Importance function

Open Issues

- Estimate other types of performance measures
- Guidance quality of state distance heuristic
- Tool implementation and complex examples
- Computational effort?



Tool TimeNET

TimeNET - **E: Waten \TimeNET-Modelle \GMPN \RESTART\SupplyPaper.xml	
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Stochastic Petri net tool – available at www.tu-ilmenau.de/sse/TimeNET

